

Unit-8: Algorithms for LTL

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Chennai Mathematical Institute

NPTEL-course

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Module 1:

**Automata-based LTL
model-checking**

Does **Transition system** satisfy **LTL formula ϕ** ?

Does **Transition system** satisfy **LTl formula** ϕ ?

Negation $\neg \phi$

Does **Transition system** satisfy **LTL formula** ϕ ?

Negation $\neg \phi$



NBA $\mathcal{A}_{\neg\phi}$

Does **Transition system** satisfy **LTL formula ϕ** ?



NBA $\mathcal{A}_{T.S}$

Negation $\neg \phi$



NBA $\mathcal{A}_{\neg \phi}$

Does **Transition system** satisfy **LTL formula ϕ** ?



NBA $\mathcal{A}_{T.S.}$

Negation $\neg \phi$



NBA $\mathcal{A}_{\neg\phi}$

Is $L(\mathcal{A}_{T.S.}) \cap L(\mathcal{A}_{\neg\phi})$ empty?

Does **Transition system** satisfy **LTL formula ϕ** ?



NBA $\mathcal{A}_{T.S.}$

Negation $\neg \phi$



NBA $\mathcal{A}_{\neg\phi}$

Is $L(\mathcal{A}_{T.S.}) \cap L(\mathcal{A}_{\neg\phi})$ empty?

Is $L(\mathcal{A}_{T.S.} \times \mathcal{A}_{\neg\phi})$ empty?

Here: Converting LTL formulas to NBA

Here: Converting LTL formulas to NBA

Coming next: Examples

Atomic propositions $\mathbf{AP} = \{ p_1, p_2 \}$

Alphabet:

$\{ \{ \}, \{ p_1 \}, \{ p_2 \}, \{ p_1, p_2 \} \}$

$\mathbb{F} p_1$ Words where p_1 occurs sometime

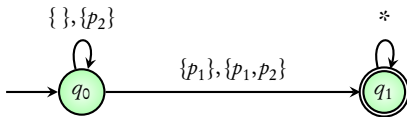
$\{p_2\} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

\vdots

F p_1 Words where p_1 occurs sometime

$\{p_2\} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$
 $\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$
 \vdots



G p_1 Words where p_1 occurs always

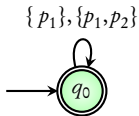
$\{p_1\} \{p_1, p_2\} \{p_1\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \dots$

$\{p_1, p_2\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

\vdots

G p_1 Words where p_1 occurs always

$\{p_1\} \{p_1, p_2\} \{p_1\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \dots$
 $\{p_1, p_2\} \{p_1, p_2\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$
 \vdots



$p_1 \wedge \neg p_2$ Words starting with $\{p_1\}$

$\{p_1\} \{\}\{\}\{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1\} \{\}\{\}\{\}\{p_1\} \{p_1\} \{p_1, p_2\} \dots$

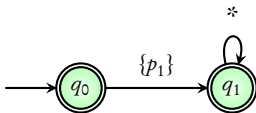
\vdots

$p_1 \wedge \neg p_2$ Words starting with $\{p_1\}$

$\{p_1\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1\} \{\} \{\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

\vdots



$$p_1 \wedge \mathbf{X} \neg p_2$$

$$\{p_1\} \{\}\{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$$

$$\{p_1, p_2\} \{p_1\} \{\}\{p_1\} \{p_1\} \{p_1, p_2\} \dots$$

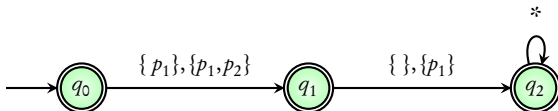
⋮

$$p_1 \wedge \mathbf{X} \neg p_2$$

$\{p_1\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1, p_2\} \{p_1\} \{\} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

\vdots

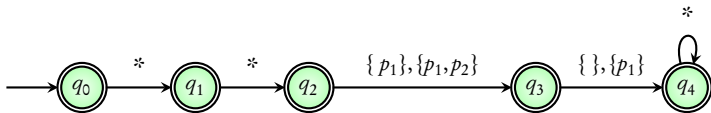


$$\mathbf{XX}(p_1 \wedge \mathbf{X}\neg p_2)$$

$$\begin{aligned} & \{ \} \{ \} \{ p_1 \} \{ \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_2 \} \{ p_2 \} \dots \\ & \{ p_2 \} \{ p_1 \} \{ p_1, p_2 \} \{ p_1 \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1, p_2 \} \dots \\ & \quad \vdots \end{aligned}$$

$$\mathbf{XX}(p_1 \wedge \mathbf{X}\neg p_2)$$

$\{\}\{\}\{p_1\}\{\}\{p_2\}\{p_1,p_2\}\{p_2\}\{p_2\}\{p_2\}\dots$
 $\{p_2\}\{p_1\}\{p_1,p_2\}\{p_1\}\{\}\{p_1\}\{p_1\}\{p_1,p_2\}\dots$
 \vdots



$$p_1 \cup p_2$$

$$\{p_1\} \{p_1\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$$

$$\{p_1, p_2\} \{ \} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$$

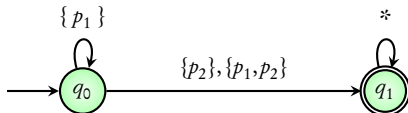
⋮

$p_1 \cup p_2$

$\{p_1\} \{p_1\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{p_1, p_2\} \{ \} \{ \} \{p_1\} \{p_1\} \{p_1, p_2\} \dots$

\vdots



$(X \ p_1) \ U \ p_2$

$\{p_2\} \{\} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{\} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \{p_1, p_2\} \dots$

$\{\} \{p_1, p_2\} \{\} \{\} \{p_2\} \{p_1, p_2\} \dots$

\vdots

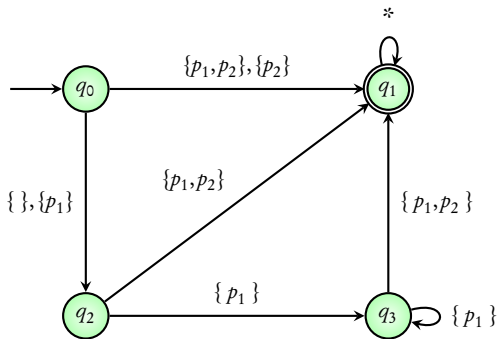
$(X p_1) U p_2$

$\{p_2\} \{ \} \{p_2\} \{p_1, p_2\} \{p_2\} \{p_2\} \{p_2\} \dots$

$\{ \} \{p_1\} \{p_1\} \{p_1\} \{p_1, p_2\} \{p_1, p_2\} \dots$

$\{ \} \{p_1, p_2\} \{ \} \{ \} \{p_2\} \{p_1, p_2\} \dots$

\vdots

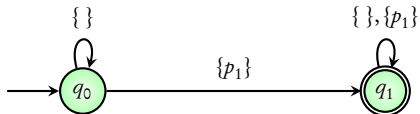


$$\mathbf{F} p_1 \wedge \neg \mathbf{F} p_2$$

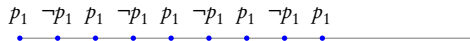
$\{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \dots$
 $\{ p_1 \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \dots$
 \vdots

$$\mathbf{F} p_1 \wedge \neg \mathbf{F} p_2$$

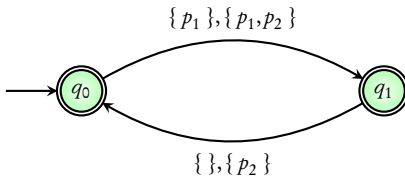
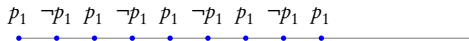
$\{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \dots$
 $\{ p_1 \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \{ \} \dots$
 \vdots



$$p_1 \wedge \mathbf{X} \neg p_1 \wedge \mathbf{G} (p_1 \leftrightarrow \mathbf{X} \mathbf{X} p_1)$$



$$p_1 \wedge \mathbf{X} \neg p_1 \wedge \mathbf{G} (p_1 \leftrightarrow \mathbf{X} \mathbf{X} p_1)$$



GF p_1 Words where p_1 occurs infinitely often

$\{ \} \{ p_1 \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_1 \} \{ p_2 \} \dots$

$\{ \} \{ \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1 \} \dots$

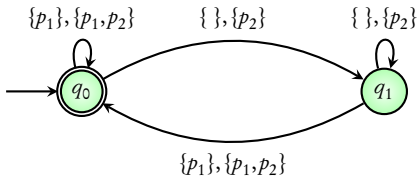
\vdots

GF p_1 Words where p_1 occurs infinitely often

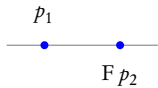
$\{ \} \{ p_1 \} \{ p_2 \} \{ p_1, p_2 \} \{ p_2 \} \{ p_1 \} \{ p_2 \} \dots$

$\{ \} \{ \} \{ \} \{ \} \{ p_1 \} \{ p_1 \} \{ p_1 \} \dots$

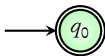
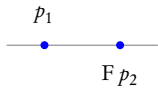
\vdots



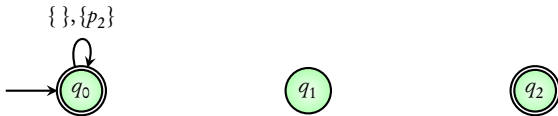
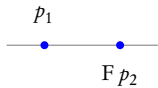
$G(p_1 \rightarrow \mathbf{XF} p_2)$



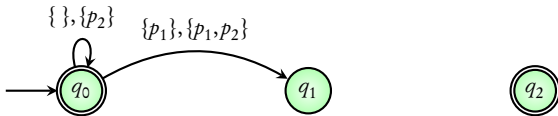
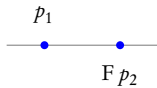
$G(p_1 \rightarrow \text{XF } p_2)$



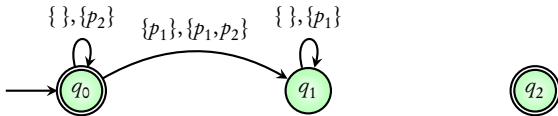
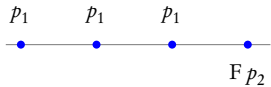
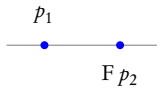
$G(p_1 \rightarrow \text{XF } p_2)$



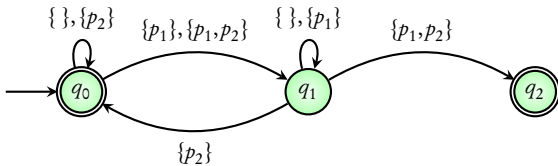
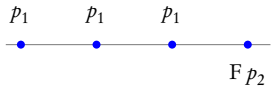
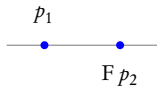
$G(p_1 \rightarrow \text{XF } p_2)$



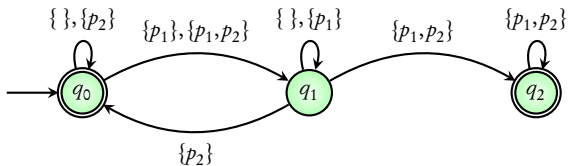
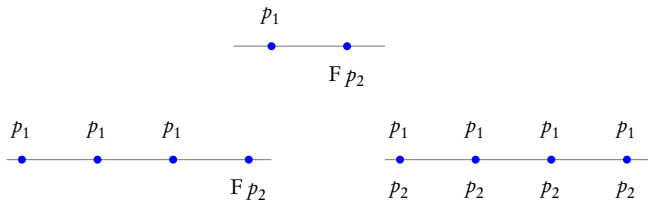
$G(p_1 \rightarrow \text{XF } p_2)$



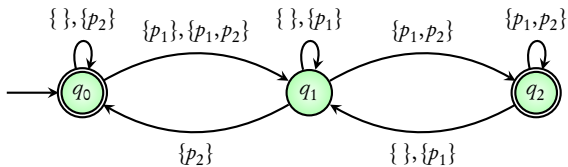
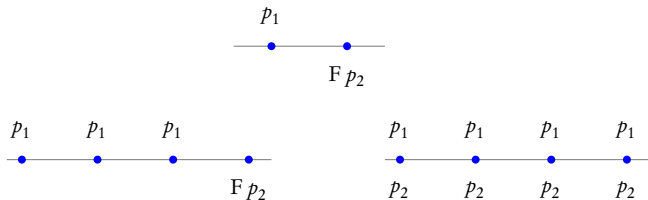
$G(p_1 \rightarrow \text{XF } p_2)$



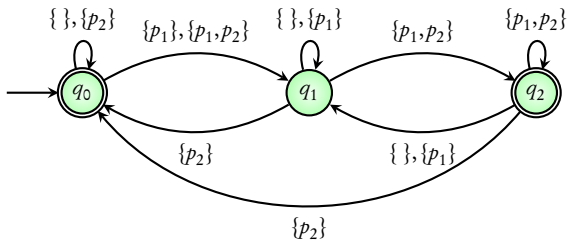
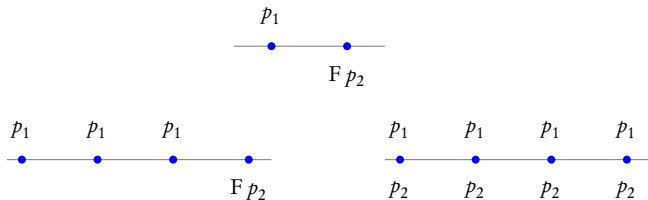
$G(p_1 \rightarrow \text{XF } p_2)$



$G(p_1 \rightarrow \text{XF } p_2)$



$G(p_1 \rightarrow \text{XF } p_2)$



Summary

LTL model-checking

Method

LTL to NBA examples

Unit-8: Algorithms for LTL

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Module 2:
LTl to NBA

Goal: Understand the **evaluation** of an LTL formula on an infinite word

$$p_1 \cup p_2$$

$$p_1 \cup p_2$$

$$\{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_2\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1\} \quad \{p_1, p_2\} \quad \dots$$

$$p_1 \cup p_2$$

$\{p_1\}$ $\{p_1\}$ $\{p_1\}$ $\{p_1\}$ $\{p_2\}$ $\{p_1\}$ $\{p_1\}$ $\{p_1\}$ $\{p_1, p_2\}$ \dots

p_1

p_2

$p_1 \cup p_2$

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$...
p_1									
p_2									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

p_1									
p_2									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
p_1									
p_2									
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
p_1	1	0	1	0	1	0	1	0	1
p_2	0	0	0	0	0	0	0	0	0
$p_1 \cup p_2$									

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\} \dots$
p_1	1	1	1	1	0	1	1	1	1
p_2	0	0	0	0	1	0	0	0	1
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\} \dots$
p_1	1	0	1	0	1	0	1	0	1
p_2	0	0	0	0	0	0	0	0	0
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0

GF p_1

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = true \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg \mathit{true} \mathbf{U} \neg \phi$

$$\neg \mathit{true} \mathbf{U} \neg(\mathit{true} \mathbf{U} p_1)$$

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = true \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$

$\{\}$ $\{\}$ $\{p_1\}$ $\{\}$ $\{\}$ $\{p_1\}$ $\{\}$ $\{\}$ $\{p_1\}$

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = \text{true} \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg \text{true} \mathbf{U} \neg \phi$

$$\neg \text{true} \mathbf{U} \neg(\text{true} \mathbf{U} p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>									
<i>true</i>									
<i>true U p₁</i>									
<i>¬ true U p₁</i>									
<i>true U ¬(true U p₁)</i>									
<i>¬ true U ¬(true U p₁)</i>									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>									
<i>true</i> \cup <i>p₁</i>									
$\neg true \cup p_1$									
<i>true</i> \cup $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> \cup <i>p₁</i>									
$\neg true \cup p_1$									
<i>true</i> \cup $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> \cup <i>p₁</i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$									
<i>true</i> \cup $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> \cup <i>p₁</i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> \cup $\neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> \cup <i>p₁</i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> \cup $\neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}	{p ₁ }
<i>p₁</i>	0	0	1	0	0	1	0	0	1
<i>true</i>	1	1	1	1	1	1	1	1	1
<i>true</i> \cup <i>p₁</i>	1	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0	0
<i>true</i> \cup $\neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1

GF p_1

GF p_1

recall that $F \phi = true \cup \phi$ and $G \phi = \neg true \cup \neg \phi$

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = true \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$

$$\mathbf{GF} p_1$$

recall that $\mathbf{F} \phi = true \mathbf{U} \phi$ and $\mathbf{G} \phi = \neg true \mathbf{U} \neg \phi$

$$\neg true \mathbf{U} \neg (true \mathbf{U} p_1)$$

{p₁} {p₁} {} {} {} {} {} {} {} {}

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1									
$true$									
$true \cup p_1$									
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$									
$true \cup p_1$									
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$									
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$									
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$									
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$									

$$GF p_1$$

recall that $F\phi = true \cup \phi$ and $G\phi = \neg true \cup \neg\phi$

$$\neg true \cup \neg(true \cup p_1)$$

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that $\mathbf{F} \phi = \text{true} \mathbf{U} \phi$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$

$$\mathit{true} \mathbf{U} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\mathbf{F} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

recall that $\mathbf{F} \phi = \mathit{true} \mathbf{U} \phi$

$$\mathit{true} \mathbf{U} (\neg p_1 \wedge \mathbf{X} (\neg p_2 \mathbf{U} p_1))$$

$$\{\} \quad \{p_2\} \quad \{\} \quad \{\} \quad \{p_1\} \quad \{p_1, p_2\} \quad \{p_1, p_2\} \quad \dots$$

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1								
p_2								
$\neg p_1$								
$\neg p_2$								
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$								
$\neg p_2$								
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$								
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$					1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p ₂ }	{}	{}	{p ₁ }	{p ₁ , p ₂ }	{p ₁ , p ₂ }	...
<i>p₁</i>	0	0	0	0	1	1	1	
<i>p₂</i>	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$			1	1	1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$								
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0							
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p ₂ }	{}	{}	{p ₁ }	{p ₁ ,p ₂ }	{p ₁ ,p ₂ }	...
<i>p</i> ₁	0	0	0	0	1	1	1	
<i>p</i> ₂	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1						
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1					
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1				
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1			
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1		
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$								
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0							
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1						
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1					
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p ₂ }	{}	{}	{p ₁ }	{p ₁ ,p ₂ }	{p ₁ ,p ₂ }	...
<i>p</i> ₁	0	0	0	0	1	1	1	
<i>p</i> ₂	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1				
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0			
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0		
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$								

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$	1	1	1	1				

$$F (\neg p_1 \wedge X (\neg p_2 U p_1))$$

recall that $F \phi = true U \phi$

$$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$$

	{}	{p ₂ }	{}	{}	{p ₁ }	{p ₁ ,p ₂ }	{p ₁ ,p ₂ }	...
<i>p</i> ₁	0	0	0	0	1	1	1	
<i>p</i> ₂	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 U p_1$	0	0	1	1	1	1	1	
$X (\neg p_2 U p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X (\neg p_2 U p_1)$	0	1	1	1	0	0	0	
$true U (\neg p_1 \wedge X (\neg p_2 U p_1))$	1	1	1	1	0	0	0	

$$p_1 \cup p_2$$

	{p ₁ }	{p ₁ }	{p ₁ }	{p ₁ }	{p ₂ }	{p ₁ }	{p ₁ }	{p ₁ }	{p ₁ ,p ₂ }	...
<i>p</i> ₁	1	1	1	1	0	1	1	1	1	
<i>p</i> ₂	0	0	0	0	1	0	0	0	0	1
<i>p</i> ₁ ∪ <i>p</i> ₂	1	1	1	1	1	1	1	1	1	1

	{p ₁ }	{}	{p ₁ }	{}	{p ₁ }	{}	{p ₁ }	{}	{p ₁ }	...
<i>p</i> ₁	1	0	1	0	1	0	1	0	1	
<i>p</i> ₂	0	0	0	0	0	0	0	0	0	
<i>p</i> ₁ ∪ <i>p</i> ₂	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	{}	{p ₂ }	{}	{}	{p ₁ }	{p ₁ ,p ₂ }	{p ₁ ,p ₂ }	...
<i>p</i> ₁	0	0	0	0	1	1	1	1
<i>p</i> ₂	0	1	0	0	0	1	1	1
¬ <i>p</i> ₁	1	1	1	1	0	0	0	0
¬ <i>p</i> ₂	1	0	1	1	1	0	0	0
¬ <i>p</i> ₂ ∪ <i>p</i> ₁	0	0	1	1	1	1	1	1
X(¬ <i>p</i> ₂ ∪ <i>p</i> ₁)	0	1	1	1	1	1	0	0
¬ <i>p</i> ₁ ∧ X(¬ <i>p</i> ₂ ∪ <i>p</i> ₁)	0	1	1	1	0	0	0	0
true ∪ (¬ <i>p</i> ₁ ∧ X(¬ <i>p</i> ₂ ∪ <i>p</i> ₁))	1	1	1	1	1	1	1	0

$$\neg true \cup \neg(true \cup p_1)$$

	{}	{}	{p ₁ }	{}	{}	{p ₁ }	{}	{}
<i>p</i> ₁	0	0	1	0	0	1	0	0
true	1	1	1	1	1	1	1	1
true ∪ <i>p</i> ₁	1	1	1	1	1	1	1	1
¬ true ∪ <i>p</i> ₁	0	0	0	0	0	0	0	0
true ∪ ¬(true ∪ <i>p</i> ₁)	0	0	0	0	0	0	0	0
¬ true ∪ ¬(true ∪ <i>p</i> ₁)	1	1	1	1	1	1	1	1

	{p ₁ }	{p ₁ }	{}	{}	{}	{}	{}	{}	{}
<i>p</i> ₁	1	1	0	0	0	0	0	0	0
true	1	1	1	1	1	1	1	1	1
true ∪ <i>p</i> ₁	1	1	0	0	0	0	0	0	0
¬ true ∪ <i>p</i> ₁	0	0	1	1	1	1	1	1	1
true ∪ ¬(true ∪ <i>p</i> ₁)	1	1	1	1	1	1	1	1	1
¬ true ∪ ¬(true ∪ <i>p</i> ₁)	0	0	0	0	0	0	0	0	0

Formula expansions

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$...
p_1	1	1	1	1	0	1	1	1	1	
p_2	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$...
p_1	1	0	1	0	1	0	1	0	1	
p_2	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
p_1	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

Key idea: Construct automata whose **states are columns of the formula expansion**

Key idea: Construct automata whose **states are columns of the formula expansion**

Next in this module: understand **properties** of formula expansions

Word compatibility

p_1								
p_2								

Word compatibility

	{ }							
p_1	0							
p_2	0							

Word compatibility

	{}		{ p_1 }					
p_1	0		1					
p_2	0		0					

Word compatibility

	{ }		{ p_1 }		{ p_2 }			
p_1	0		1		0			
p_2	0		0		1			

Word compatibility

	$\{\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1, p_2\}$
p_1	0	1	0	1
p_2	0	0	1	1

AND-NOT-compatibility

ϕ

	0		1	
--	---	--	---	--

$\neg\phi$

	1		0	
--	---	--	---	--

AND-NOT-compatibility

 ϕ

	0		1	
--	---	--	---	--

 $\neg\phi$

	1		0	
--	---	--	---	--

 ϕ_1

1		0		1		0
---	--	---	--	---	--	---

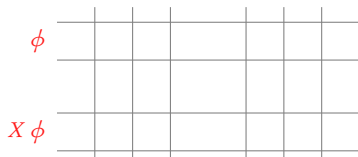
 ϕ_2

1		1		0		0
---	--	---	--	---	--	---

 $\phi_1 \wedge \phi_2$

1		0		0		0
---	--	---	--	---	--	---

X-compatibility



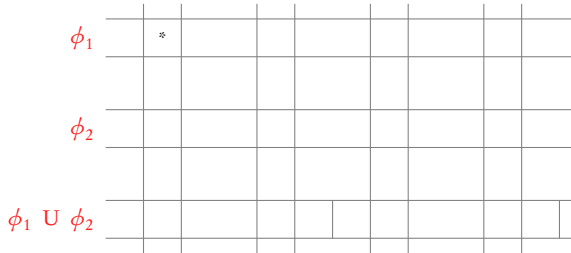
X-compatibility

ϕ		0				
$X\phi$	0					

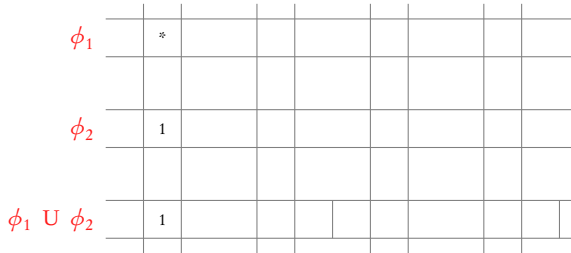
X-compatibility

ϕ		0			1	
$X\phi$	0				1	

Until-compatibility



Until-compatibility



Until-compatibility

ϕ_1	*						
ϕ_2	1		0				
$\phi_1 \text{ U } \phi_2$	1		1				

Until-compatibility

ϕ_1	*	1					
ϕ_2	1	0					
$\phi_1 \text{ U } \phi_2$	1	1	1				

Until-compatibility

ϕ_1	*		1					
ϕ_2	1		0		0			
$\phi_1 \text{ U } \phi_2$	1		1	1		0		

Until-compatibility

ϕ_1	*		1		0			
ϕ_2	1		0		0			
$\phi_1 \text{ U } \phi_2$	1		1	1		0		

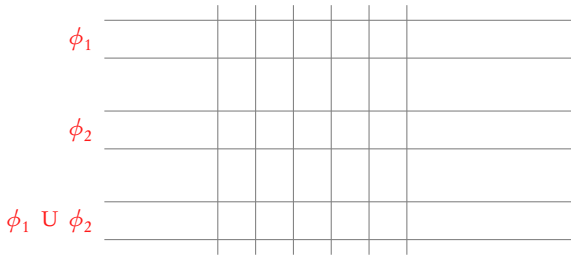
Until-compatibility

ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \text{ U } \phi_2$	1		1	1		0		0

Until-compatibility

ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \text{ U } \phi_2$	1		1	1		0		0

Until-compatibility: eventuality condition



Until-compatibility: eventuality condition

ϕ_1	1					
ϕ_2	0					
$\phi_1 \text{ U } \phi_2$	1	1				

Until-compatibility: eventuality condition

ϕ_1		1	1				
ϕ_2		0	0				
$\phi_1 \text{ U } \phi_2$		1	1	1			

Until-compatibility: eventuality condition

ϕ_1		1	1	1			
ϕ_2		0	0	0			
$\phi_1 \text{ U } \phi_2$		1	1	1	1		

Until-compatibility: eventuality condition

ϕ_1		1	1	1	1		
ϕ_2		0	0	0	0		
$\phi_1 \text{ U } \phi_2$		1	1	1	1	1	

Until-compatibility: eventuality condition

ϕ_1		1	1	1	1	1	
ϕ_2		0	0	0	0	0	...
$\phi_1 \text{ U } \phi_2$		1	1	1	1	1	

Until-compatibility: eventuality condition

ϕ_1		1	1	1	1	1	
ϕ_2		0	0	0	0	0	...
$\phi_1 \cup \phi_2$		1	1	1	1	1	

Cannot happen forever that $\phi_1 \cup \phi_2 = 1$, $\phi_1 = 1$ but $\phi_2 = 0$

Accepting expansions

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_2\}$	$\{p_1, p_2\}$	$\{\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1, p_2\}$	$\{\}$	$\{p_1, p_2\}$	\dots
p_1	1	1	1	1	0	1	1	1	1	
p_2	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_2\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_2\}$	$\{\}$	$\{p_1\}$	\dots
p_1	1	0	1	0	1	0	1	0	1	
p_2	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$\text{true} \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$	\dots
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$\text{true} \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg \text{true} \cup \neg(\text{true} \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$
p_1	0	0	1	0	0	1	0	0
true	1	1	1	1	1	1	1	1
$\text{true} \cup p_1$	1	1	1	1	1	1	1	1
$\neg \text{true} \cup p_1$	0	0	0	0	0	0	0	0
$\text{true} \cup \neg(\text{true} \cup p_1)$	0	0	0	0	0	0	0	0
$\neg \text{true} \cup \neg(\text{true} \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
true	1	1	1	1	1	1	1	1	1
$\text{true} \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg \text{true} \cup p_1$	0	0	1	1	1	1	1	1	1
$\text{true} \cup \neg(\text{true} \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg \text{true} \cup \neg(\text{true} \cup p_1)$	0	0	0	0	0	0	0	0	0

Entry in **first column** of **last row** (corresponding to final formula) is 1

Summary

LTL to NBA

Formula expansions

Properties

Columns as states of NBA

Unit-8: Algorithms for LTL

B. Srivathsan

Chennai Mathematical Institute

NPTEL-course

July - November 2015

Module 3:
Automaton construction

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$...
p_1	1	1	1	1	0	1	1	1	1	
p_2	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$...
p_1	1	0	1	0	1	0	1	0	1	
p_2	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
p_1	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

$$p_1 \cup p_2$$

	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_2\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1\}$	$\{p_1, p_2\}$...
p_1	1	1	1	1	0	1	1	1	1	
p_2	0	0	0	0	1	0	0	0	1	
$p_1 \cup p_2$	1	1	1	1	1	1	1	1	1	

	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{p_1\}$...
p_1	1	0	1	0	1	0	1	0	1	
p_2	0	0	0	0	0	0	0	0	0	
$p_1 \cup p_2$	0	0	0	0	0	0	0	0	0	

$$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$$

	$\{\}$	$\{p_2\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{p_1, p_2\}$	$\{p_1, p_2\}$...
p_1	0	0	0	0	1	1	1	
p_2	0	1	0	0	0	1	1	
$\neg p_1$	1	1	1	1	0	0	0	
$\neg p_2$	1	0	1	1	1	0	0	
$\neg p_2 \cup p_1$	0	0	1	1	1	1	1	
$X(\neg p_2 \cup p_1)$	0	1	1	1	1	1	0	
$\neg p_1 \wedge X(\neg p_2 \cup p_1)$	0	1	1	1	0	0	0	
$true \cup (\neg p_1 \wedge X(\neg p_2 \cup p_1))$	1	1	1	1	1	1	0	

$$\neg true \cup \neg(true \cup p_1)$$

	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{p_1\}$	$\{\}$	$\{\}$
p_1	0	0	1	0	0	1	0	0
$true$	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	1	1	1	1	1	1
$\neg true \cup p_1$	0	0	0	0	0	0	0	0
$true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0
$\neg true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1

	$\{p_1\}$	$\{p_1\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$	$\{\}$
p_1	1	1	0	0	0	0	0	0	0
$true$	1	1	1	1	1	1	1	1	1
$true \cup p_1$	1	1	0	0	0	0	0	0	0
$\neg true \cup p_1$	0	0	1	1	1	1	1	1	1
$true \cup \neg(true \cup p_1)$	1	1	1	1	1	1	1	1	1
$\neg true \cup \neg(true \cup p_1)$	0	0	0	0	0	0	0	0	0

Construct an automaton with states as column vectors that can guess accepting expansions

Example 1: $p_1 \cup p_2$

p_1	0	0	0	0	1	1	1	1
p_2	0	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	1	0	1	0	1	0	1

p_1	0	0	0	0	1	1	1	1
p_2	0	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	1	0	1	0	1	0	1

ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

Recall Until-compatibility

p_1	0	0	0	1	1	1	1
p_2	0	1	1	0	0	1	1
$p_1 \cup p_2$	0	0	1	0	1	0	1

ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

Recall Until-compatibility

p_1	0	0	1	1	1	1
p_2	0	1	0	0	1	1
$p_1 \cup p_2$	0	1	0	1	0	1

ϕ_1	*	1	0	1	
ϕ_2	1	0	0	0	
$\phi_1 \cup \phi_2$	1	1	1	0	0

Recall Until-compatibility

p_1	0	0	1	1	1
p_2	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

ϕ_1	*	1	0	1		
ϕ_2	1	0	0	0		
$\phi_1 \cup \phi_2$	1	1	1	0	0	0

Recall Until-compatibility

p_1	0	0	1	1	1
p_2	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

Compatible states

ϕ_1	*	1	0	1		
ϕ_2	1	0	0	0		
$\phi_1 \cup \phi_2$	1	1	1	0	0	0

Recall Until-compatibility

p_1	0	0	1	1	1
p_2	0	1	0	0	1
$p_1 \cup p_2$	0	1	0	1	1

q_0

p_1	0
p_2	0
$p_1 \cup p_2$	0

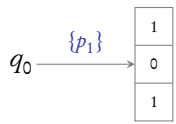
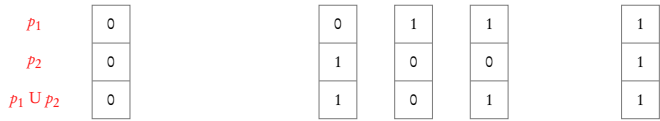
0
1
1

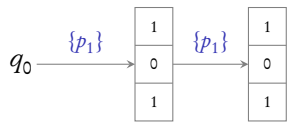
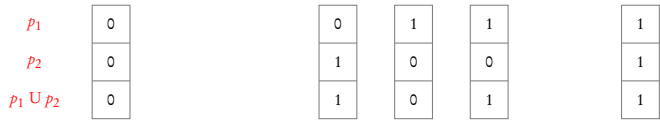
1
0
0

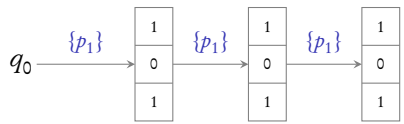
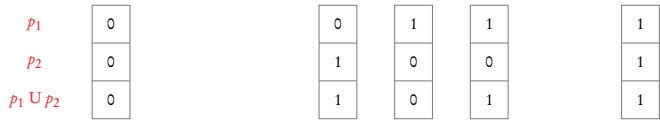
1
0
1

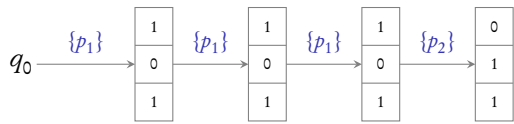
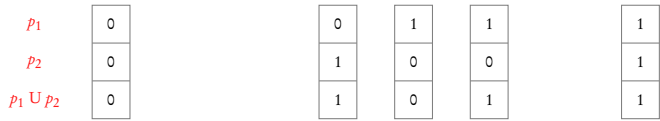
1
1
1

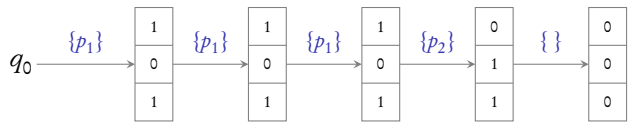
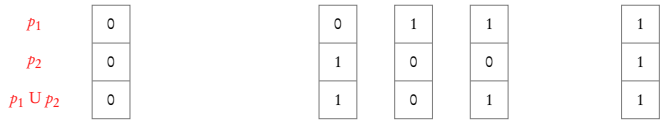
$q_0 \xrightarrow{\{p_1\}}$

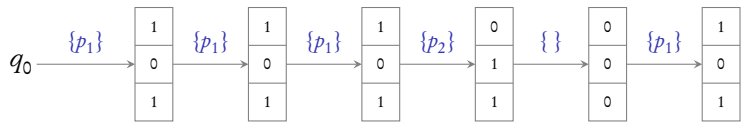
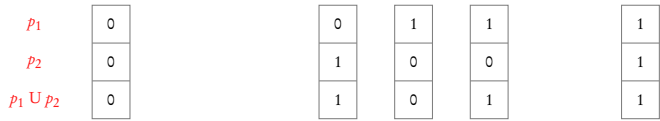


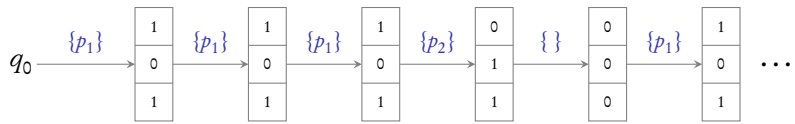
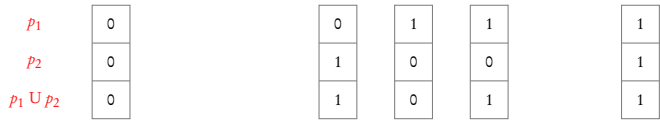


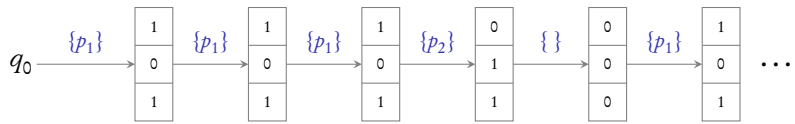
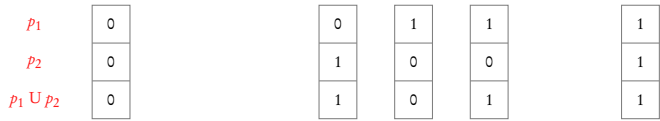




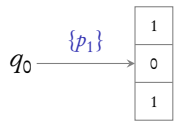
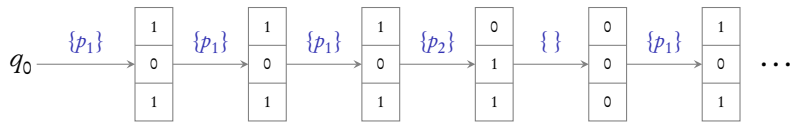
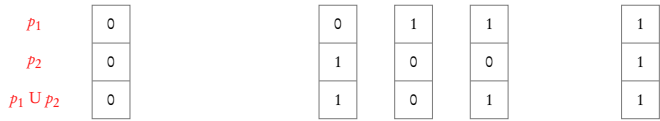


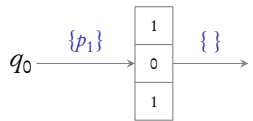
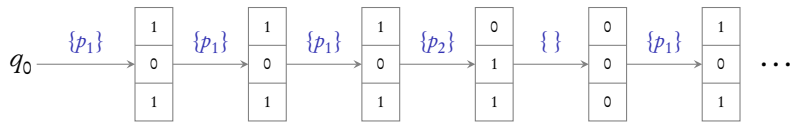
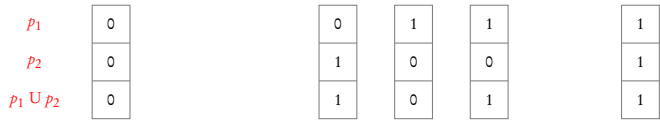


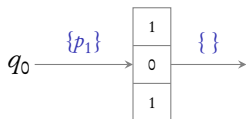
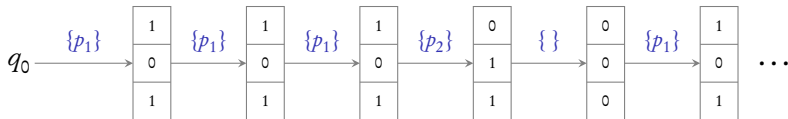
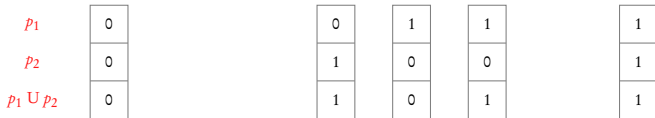




q_0







No compatible transition

ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

1
0
1

$\longrightarrow q_0$

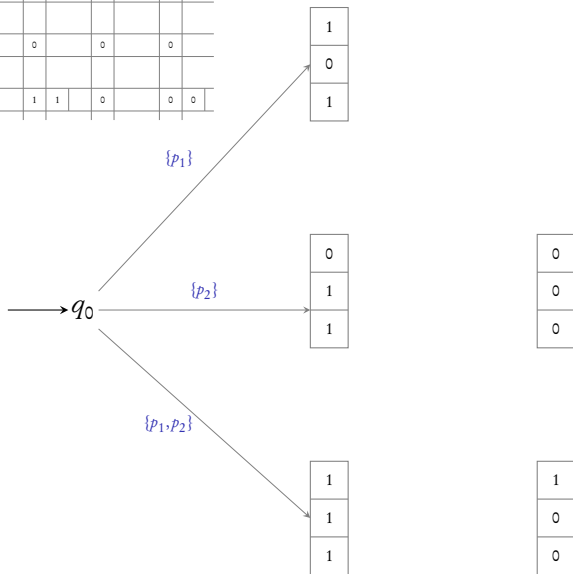
0
1
1

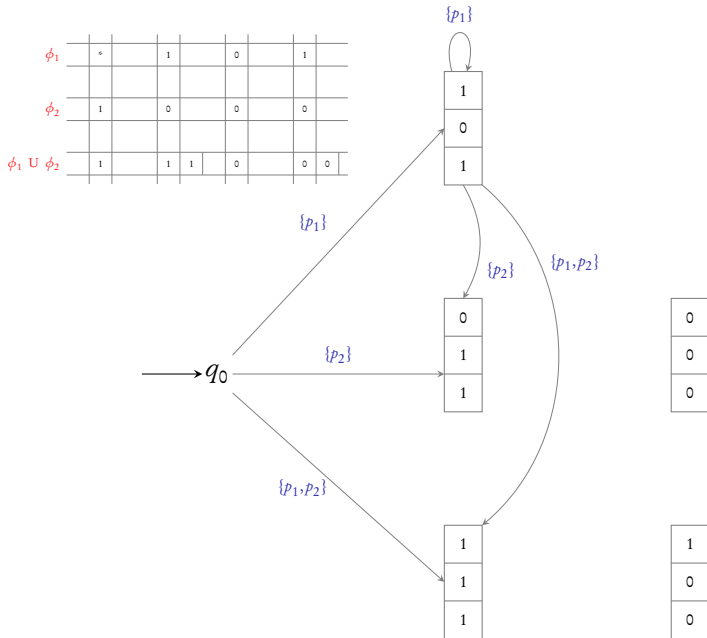
0
0
0

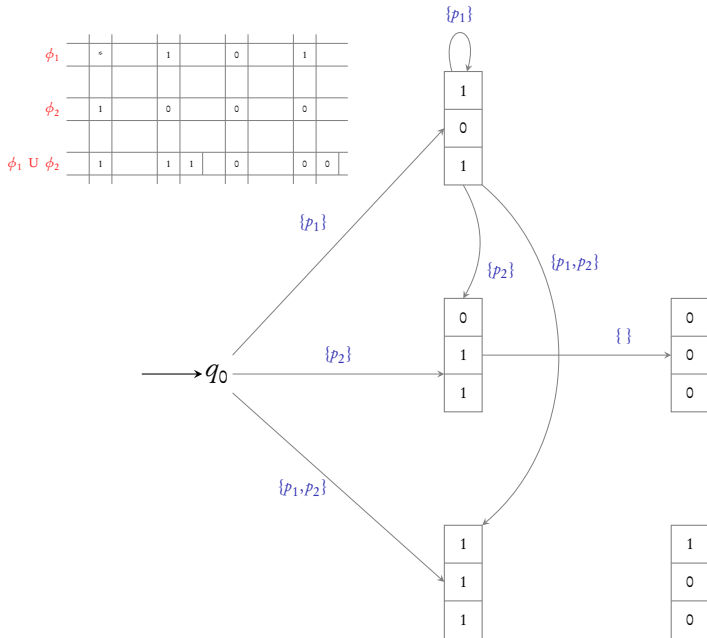
1
1
1

1
0
0

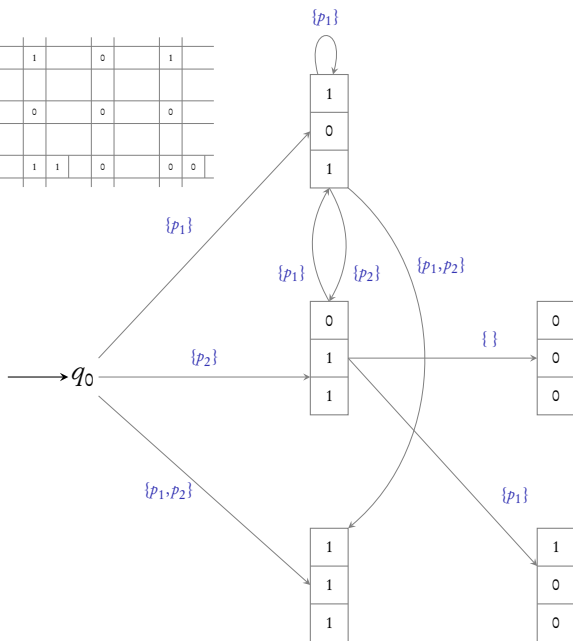
ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



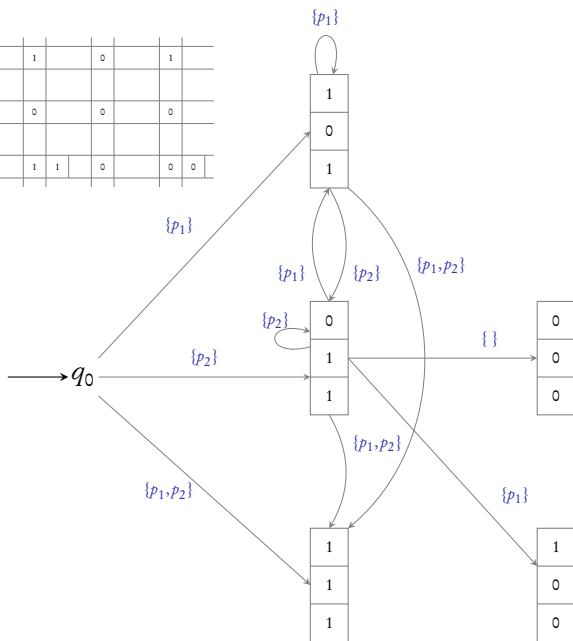




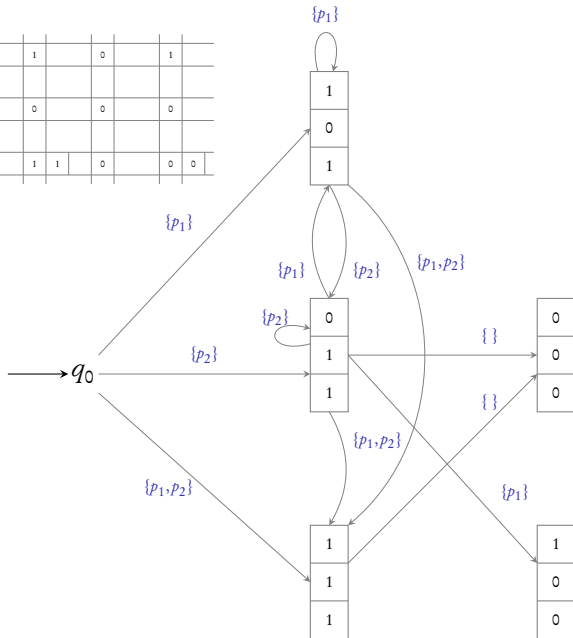
ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



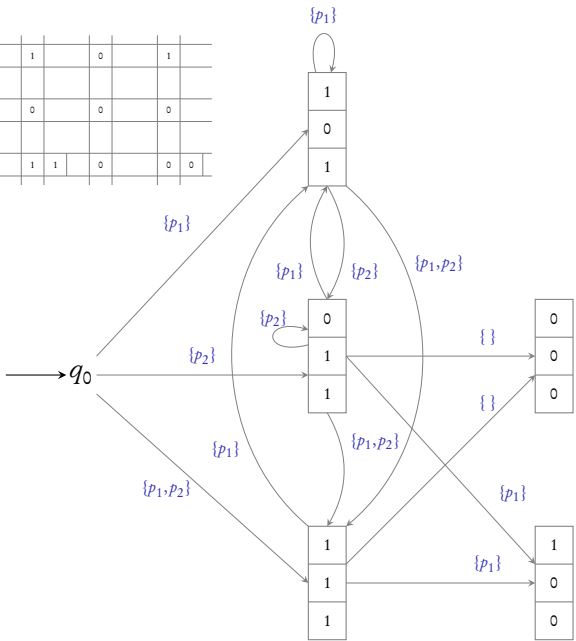
ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



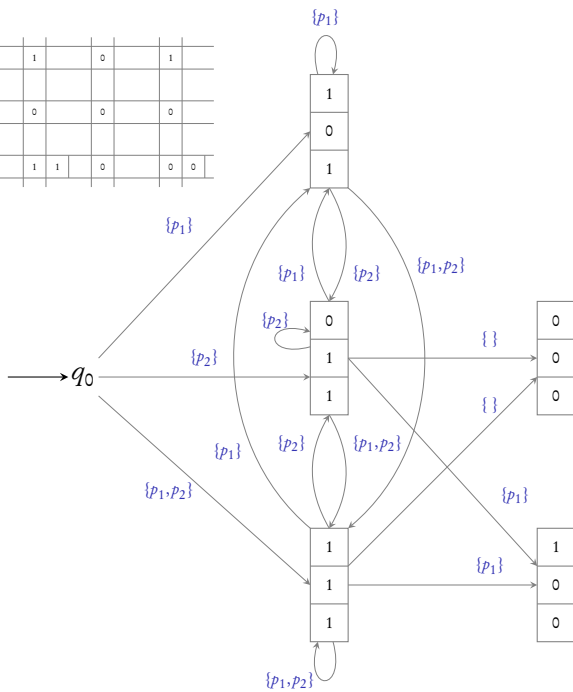
ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



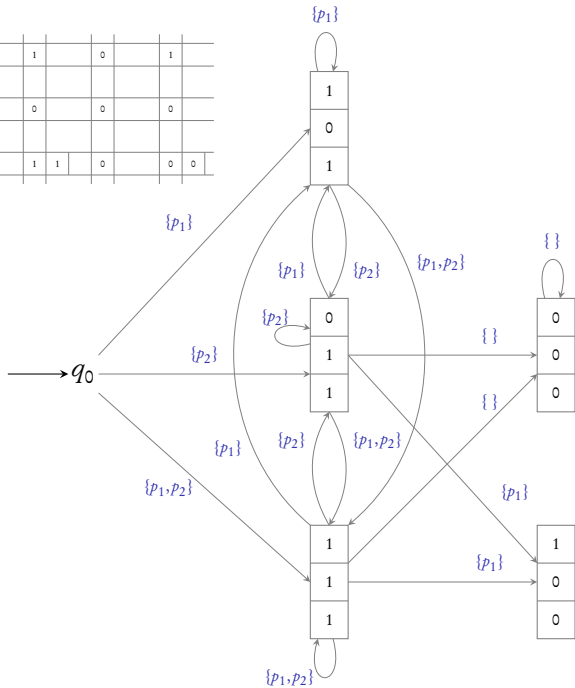
ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



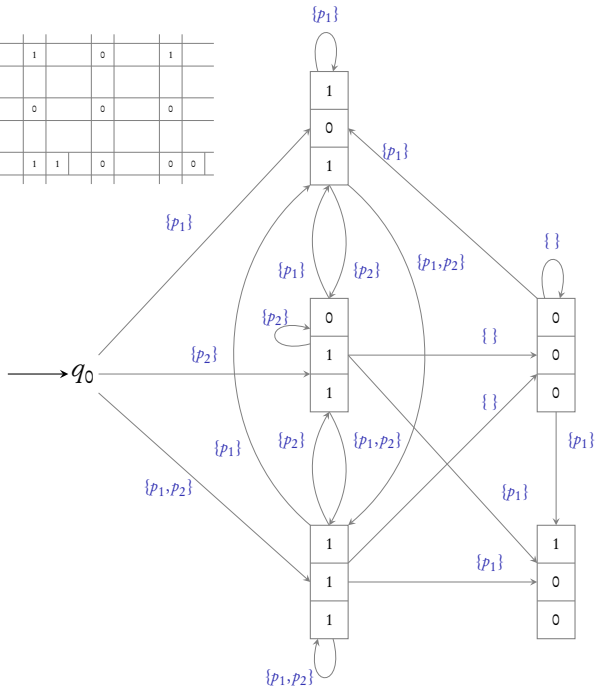
ϕ_1	*		1		0		1	
ϕ_2			0		0		0	
$\phi_1 \cup \phi_2$		1		1	1		0	0



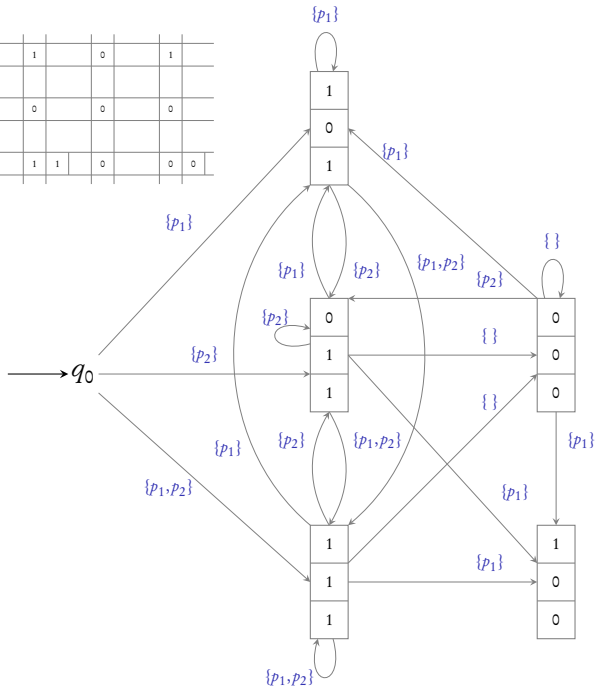
ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



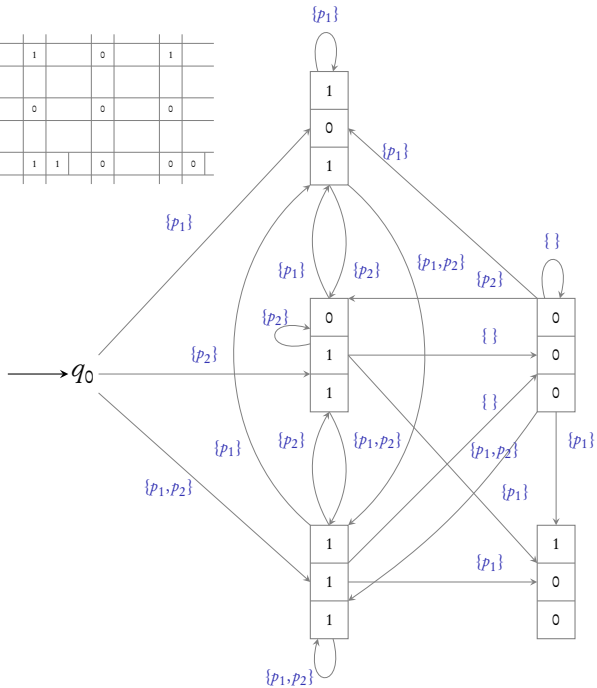
ϕ_1	*		1		0		1	
ϕ_2			0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0



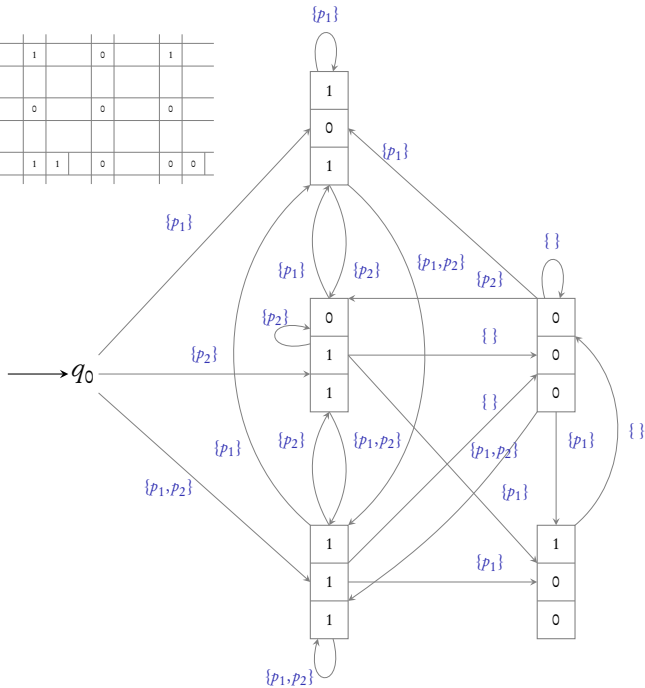
ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



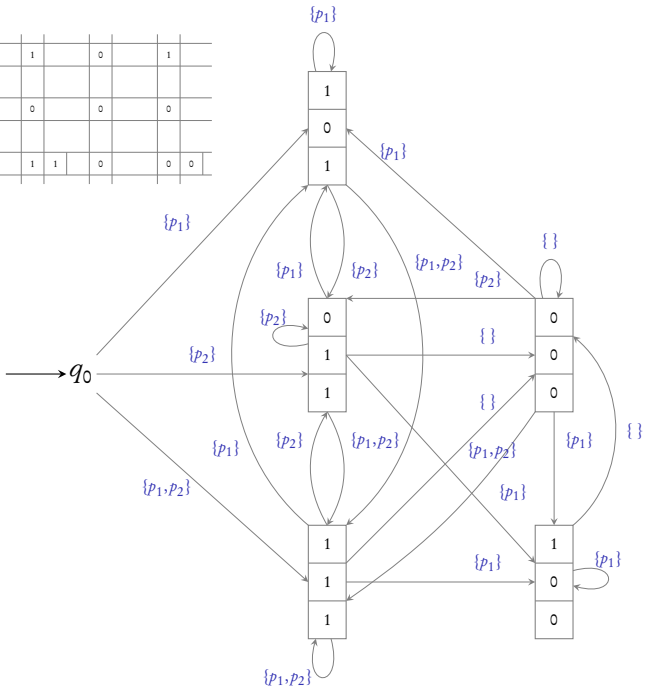
ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



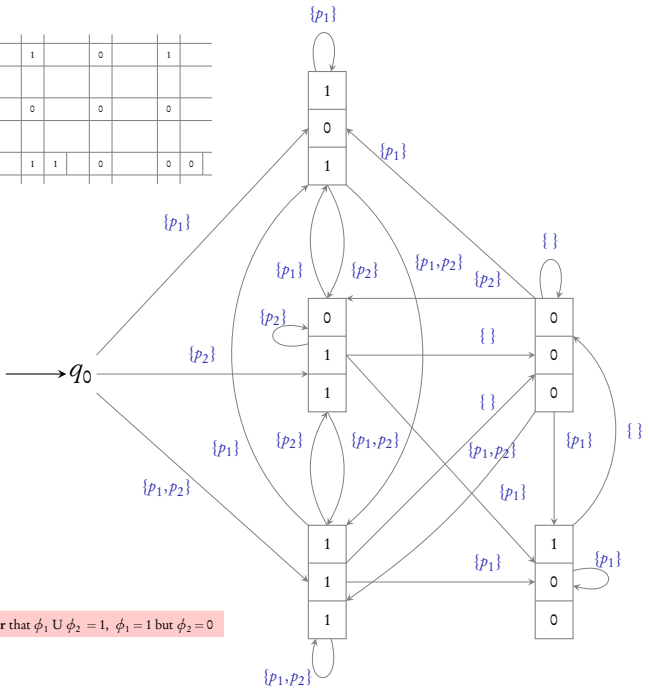
ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



ϕ_1	*		1		0		1	
ϕ_2	1		0		0		0	
$\phi_1 \cup \phi_2$	1		1	1	0		0	0

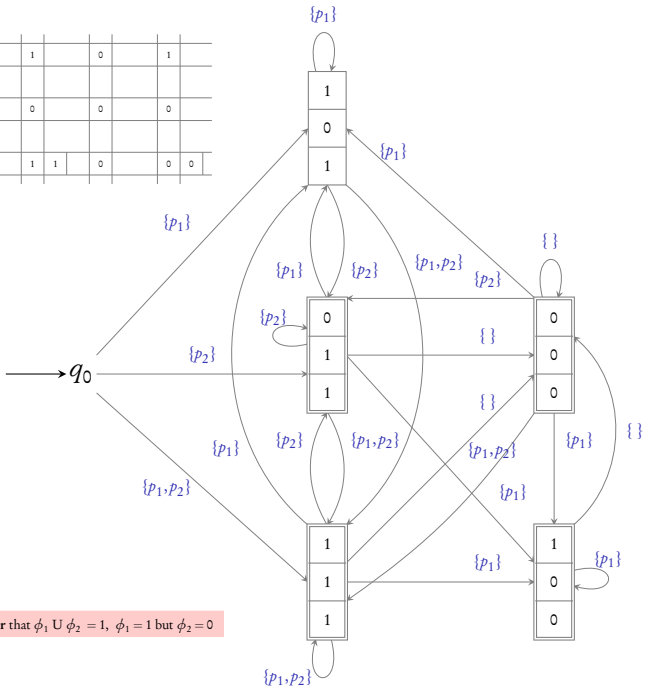


ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



Cannot happen forever that $\phi_1 \cup \phi_2 = 1$, $\phi_1 = 1$ but $\phi_2 = 0$

ϕ_1	*		1		0		1	
ϕ_2		1		0		0		0
$\phi_1 \cup \phi_2$		1		1	1		0	0



Cannot happen forever that $\phi_1 \cup \phi_2 = 1$, $\phi_1 = 1$ but $\phi_2 = 0$

Example 2: $(X p_1) \cup p_2$

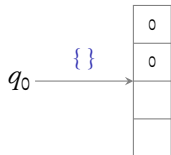
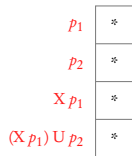
p_1	*
p_2	*
$X p_1$	*
$(X p_1) \cup p_2$	*

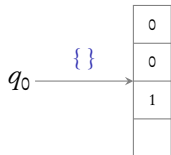
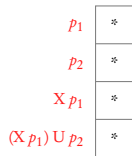
p_1	*
p_2	*
$X p_1$	*
$(X p_1) \cup p_2$	*

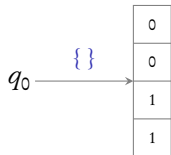
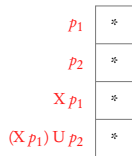
q_0

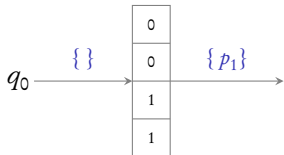
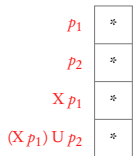
p_1	*
p_2	*
$X p_1$	*
$(X p_1) \cup p_2$	*

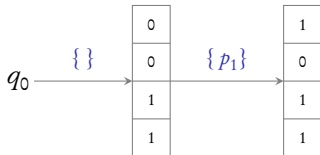
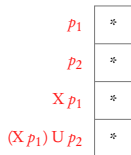
$q_0 \longrightarrow \{\}$

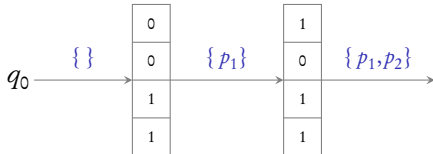
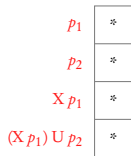


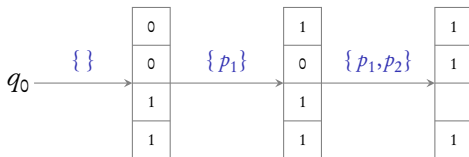
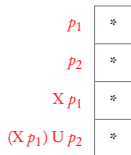


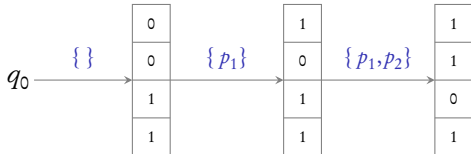
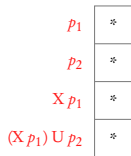


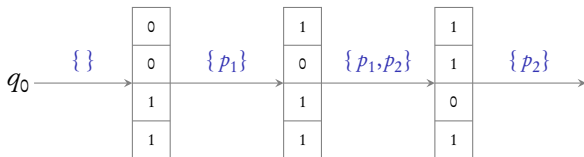
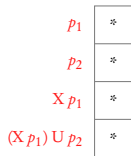


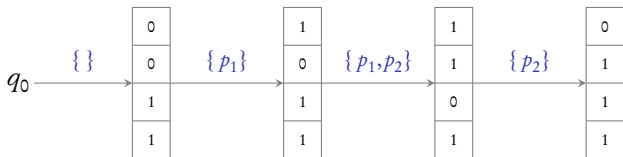
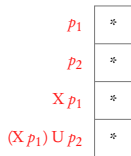


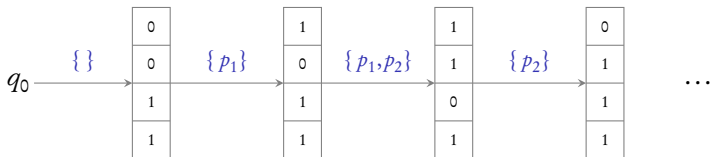
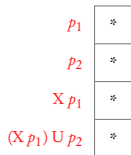












Coming next: Construction for an arbitrary LTL formula ϕ

Step 1: List down subformulae of ϕ

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p_1	*
p_2	*
$p_1 \cup p_2$	*

p_1	*
p_2	*
$\neg p_1$	*
$(\neg p_1) \cup p_2$	*

Step 2: Check **AND-NOT** and **Until** compatibility

Step 2: Check **AND-NOT** and **Until** compatibility

p_1	0
p_2	0
$p_1 \cup p_2$	1

p_1	0
p_2	1
$\neg p_1$	0
$(\neg p_1) \cup p_2$	0

Incompatible states!

Step 2: Check **AND-NOT** and **Until** compatibility

p_1	0
p_2	0
$p_1 \cup p_2$	1

p_1	0
p_2	1
$\neg p_1$	0
$(\neg p_1) \cup p_2$	0

Incompatible states!

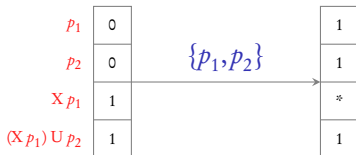
Remove incompatible states and **add** a new state $\{q_0\}$

Step 3: Add transitions satisfying

Word, X and Until compatibility

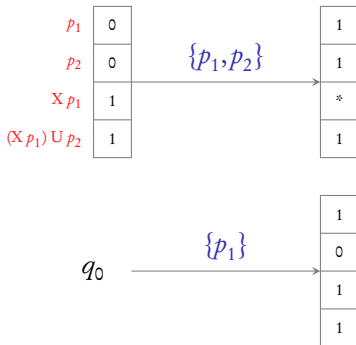
Step 3: Add transitions satisfying

Word, X and Until compatibility



Step 3: Add transitions satisfying

Word, X and Until compatibility



From q_0 add compatible transitions to states where last entry is 1

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0	1	0	1
0	0	1	1
0	0	1	1

Final automaton \mathcal{A}_ϕ

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In general, this algorithm gives NBA which is **exponential** in size of formula